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Cross-Layer Optimisation to Maximise Fairness among TCP flows of different TCP flavors.

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Presentation Structure

- Introduce the challenges and locate this research.
- TCP flavors and Packet Losses
 - **Throughput Comparison**
- Definition of Fairness
- Jain's Index as a notion of Fairness
- Improving Fairness among TCP flows
 - **Fairness Maximization Problem**
 - **Real-time solution**
- Performance Investigation
 - **Significant increase in Fairness**
 - **Minimal effect on the aggregated throughput of TCP**
- Conclusions



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Introduce the challenges

- In wired networks, end-to-end data transmission relies on the functionalities of the Transport Layer, i.e., TCP, for congestion control, fairness, loss recovery, etc.
- *Random* and *Time-variant* characteristics of wireless links fool TCP algorithms.
- A vast amount of research has been conducted in order to introduce enhancements in the Transport Layer to perform better in the wireless networks.
- A common approach is to increase the intelligence of wireless links to be aware of TCP.



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Locate this Research

- We expand the cross-layer algorithms from the traditional Layer 1 and 2 interactions, up to the Transport Layer.
- As we aim to support higher data rates, Transport layer performance plays an extremely important role.
- Our research does not attempt to propose any new technique to the Transport Layer.
- We aim to:
 - improve the end-to-end performance, which is all that can be actually seen from user's perspective.
 - rely on the well-used Transport Layer over Internet, TCP, for data transmission.
 - study a scenario that different TCP flavors competing over the wireless network.



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TCP Flavors and Packet losses

- In the absence of losses TCP congestion control behavior is similar in most TCP flavors; per RTT *cwnd* increases.
 - Exponential Initially: slow-start phase
 - Linear afterward: congestion control phase
- In the presence of losses, different TCP flavors react differently to the loss.
- TCP Throughput models, which captures this effect, are used from the literature [1], [2], and [3] for the Reno, NewReno and westwood TCP.



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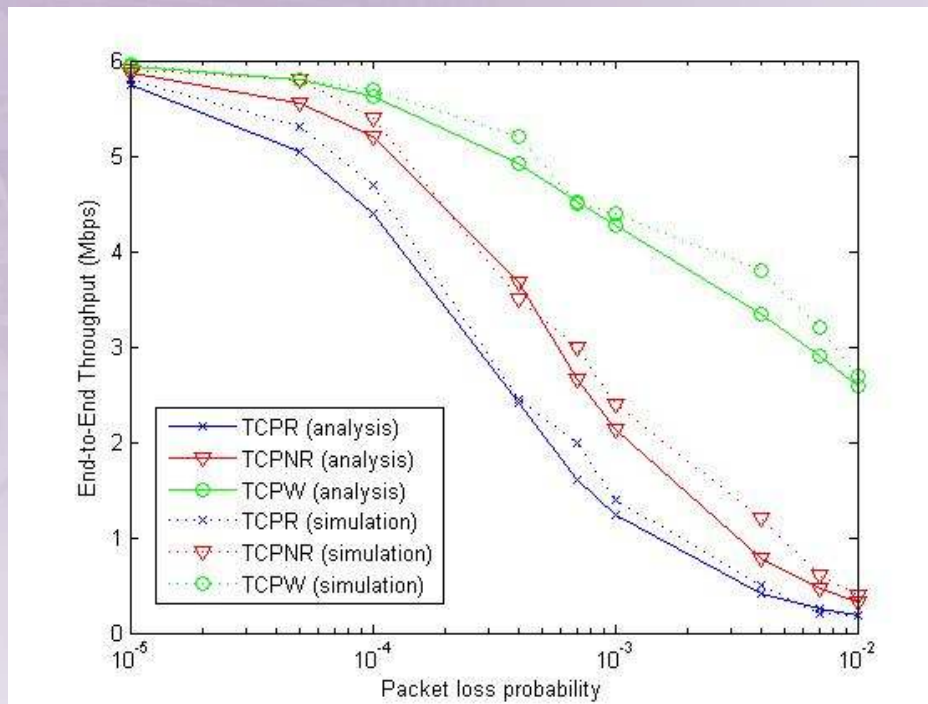


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Throughput Comparison

- Single-flow Scenario of TCP Reno, TCP NewReno, TCP westwood and SACK option.



- ✓ 6Mbps bottleneck link.
- ✓ RTT = 100ms.
- ✓ FTP download a file of ~11,000 packets
- ✓ TCP packet size is 1460B.

- TCPR degrades most dramatically in the event of losses.
- models and the analyses perform almost identically.

Jain's Fairness Index

- Jain's Index introduced to measure fairness in computer networks in 1984 [5].

$$J = \frac{\left(\sum_{i=1}^n x(p_i) \right)^2}{n \cdot \sum_{i=1}^n x(p_i)^2}$$

- ✓ J is independent of the scale of allocation metric (x).
- ✓ J is bounded between 0 and 1.
- ✓ 1 is completely fair and 0 is completely unfair allocation.

- Assuming allocated rate for user i , operating in PER p_i , is $r(p_i)$ where the optimal rate for user i is R_i , then $x(p_i) = r(p_i) / R_i$.
- n is the number of flows in the network.



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Fairness Definition

- In the literature, *Fairness* is defined as Equality.
- We define *Fairness* as Equal access to the resources proportional to the potential of each user to utilize that resources.
- For example, if users' capacity is defined by their buffer size, then the fair allocation would be to allocate more to the user with the larger buffer.

User 1

10
9
8
7
6
5
4
3
2
1

User 2

2
1

User 3

5
4
3
2
1

- ✓ Total available data is 9B.
- ✓ Equal assignment results in Jain index = 0.7.
- ✓ Assigning proportional to the users' capacity (5B to U1, 1B to U2, and 3B to U3) results in Jain index = 0.99.



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Fairness Maximisation

- The aim is to maximise fairness among TCP flows.

$$(P) : \text{ maximise } J(p) = \frac{\left(\sum_{i=1}^n x(p_i) \right)^2}{n \cdot \sum_{i=1}^n x(p_i)^2},$$

$$\sum_{i=1}^n B_k(p_i) \leq W, \quad \forall k \in \{1..4\}$$
$$p_i \cdot 10^{-e} \leq p_i \leq p_i \cdot 10^e, \quad \forall i \in \{1..n\}$$
$$p_i \geq 0, \quad \forall i \in \{1..n\}$$

- ✓ PER can be seen by flow i is p_i ; thus flow i of TCP flavor k can achieve throughput of $B_k(p_i)$.
- ✓ Throughput expressions of the four used TCPs are from [2-5]
- ✓ Optimal throughput for flow i is achieved when all the other $n-1$ users have the same TCP flavor.

$$x(p_i) = \frac{B_k(p_i)}{B_{k_{Optimal}}},$$

- ✓ Link Capacity is limited to W .
- ✓ Exponent of PER can vary in the close interval of $[p-e, p+e]$.



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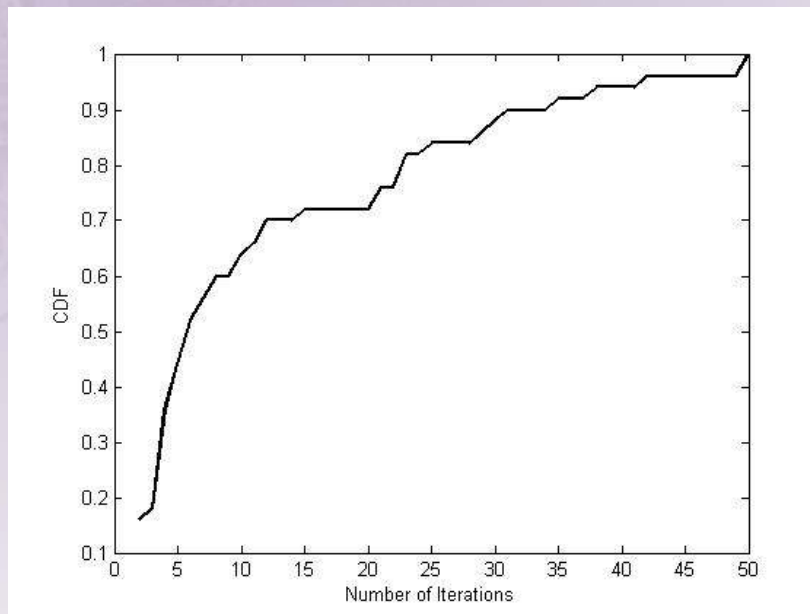


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Heuristic Solution

- The problem is approximated, to be solved it in *real time*. The approximated problem is solved using *Newton method*.
- We show that the solution converges to the optimal value in few number of iterations.



CDF of the number of iterations required in solving the problem.

- ✓ Maximum 50 iterations are used.
- ✓ Average number of required iterations is 12.
- ✓ In 70% of cases, the optimal value is attained in less than 15 iterations.



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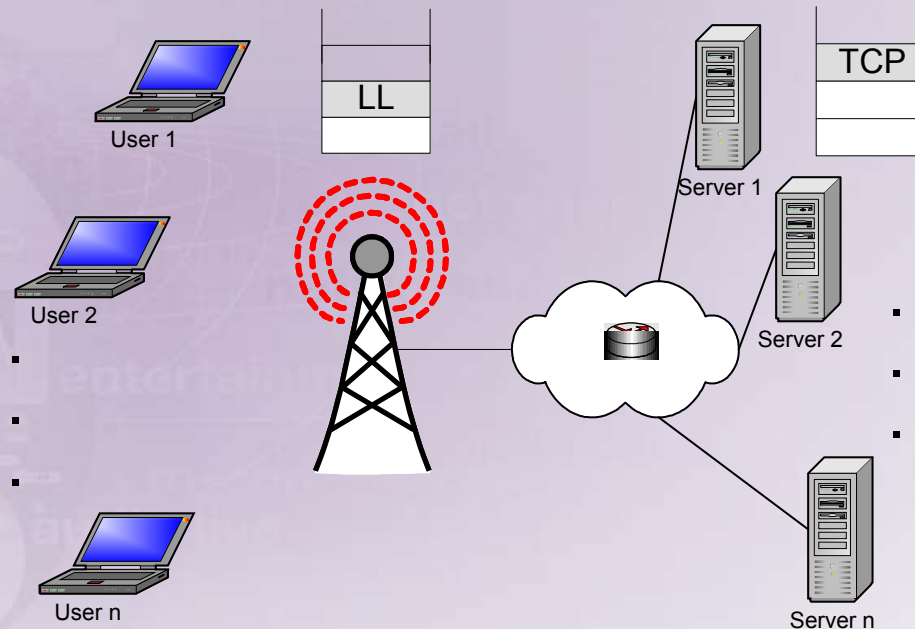
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Network Configuration

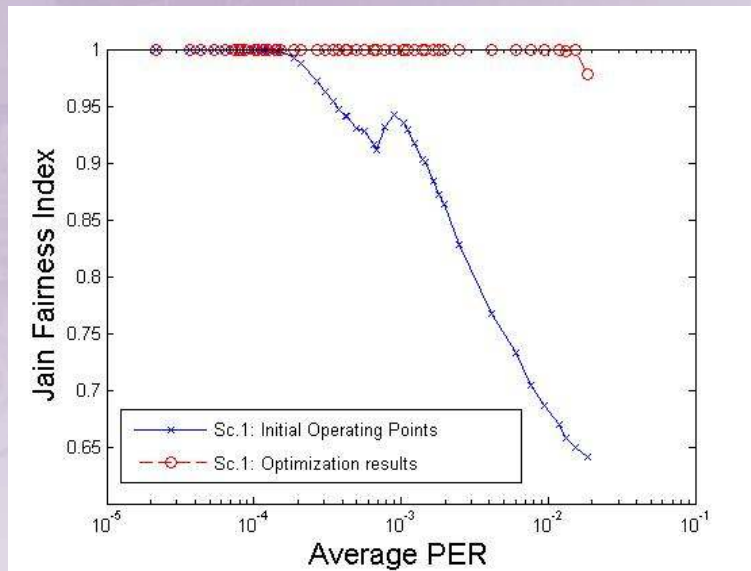


- ✓ IEEE 802.11a.
- ✓ Max data rate of 54 Mbps.
- ✓ 15 mobile users.
- ✓ 15 FTP servers.
- ✓ Each client is connected to a unique server.
- ✓ All traffic flows are served by the single wireless AP.
- ✓ TCP flows are of different combinations of the three TCP flavors.

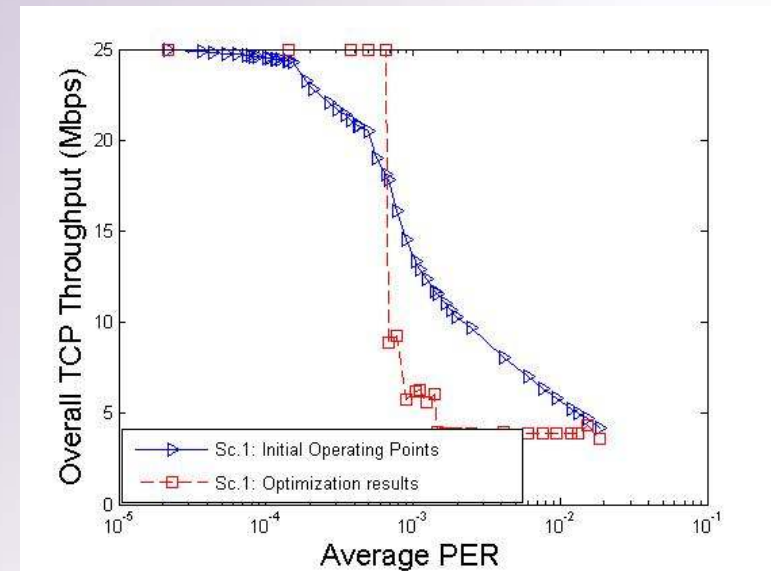
Using the information of the Transport layer, optimisation performs in the wireless AP.

Numerical Results: Scenario 1

- Of the 15 TCP flows, there are five TCPR, five TCPNR, five TCPW.



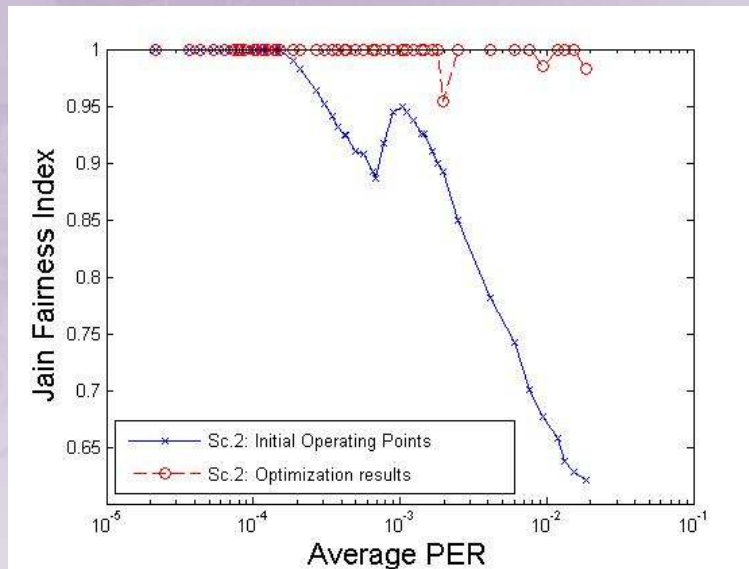
The fairness index versus average PER:
Significant increase of up to 50% in the high error rate conditions is observed.



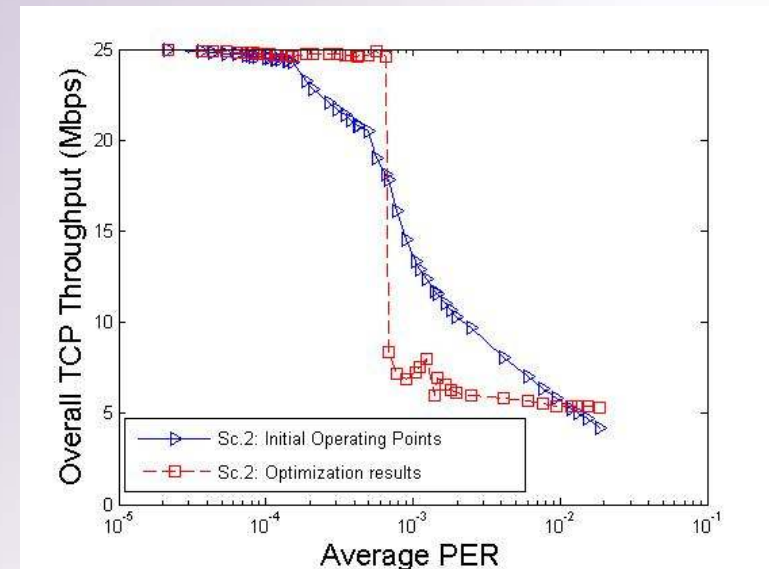
TCP throughput versus average PER:
Slight decrease in the high error rate and slight increased in the low error rate is observed.

Numerical Results: Scenario 2

- Of the 15 TCP flows, there are nine TCPR, four TCPNR, two TCPW.



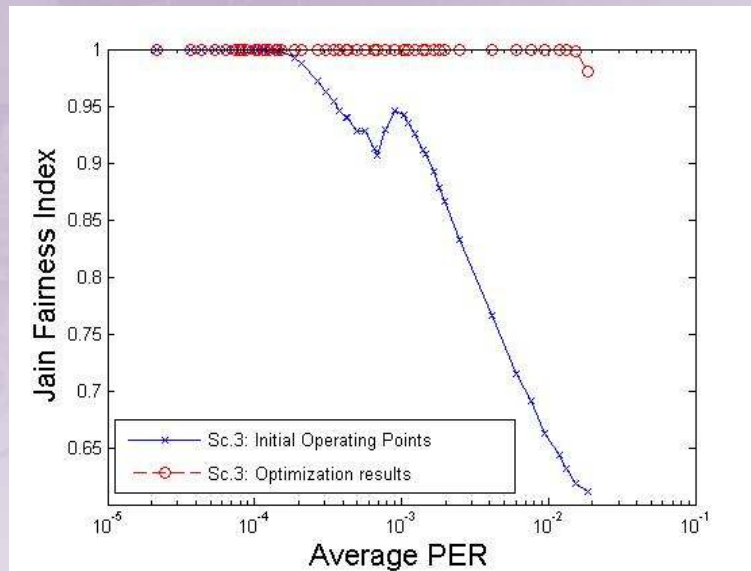
The fairness index versus average PER:
Observation is similar to Scenario 1.



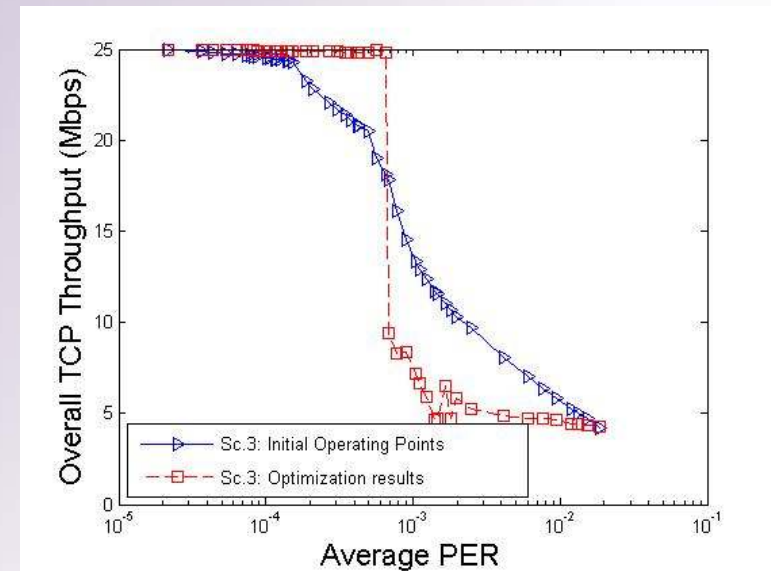
TCP throughput versus average PER:
Observation is similar to Scenario 1.

Numerical Results: Scenario 3

- Of the 15 TCP flows, there are three TCPR, eight TCPNR, four TCPW.



The fairness index versus average PER:
Observation is similar to Scenario 1.



TCP throughput versus average PER:
Observation is similar to Scenario 1.

Conclusions

- We have studied scenarios where multiple TCP flows are from different TCP flavors.
- The unfairness among TCP flows arises from different reaction to packet loss of each TCP flavor.
- The FEC code rate is adopted not only to achieve higher reliability on the wireless link, but also to consider the theoretical bands apply by the TCP flavor in use.
- Our approach doesn't attempt to change the TCP.
- The overall results show up to 50% improvement in the fairness index with no significant effect on the overall end-to-end throughput.



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