

# Enumerating words with forbidden factors

Some things that could be better known

Richard Pinch

Institute of Mathematics and its Applications

07 February 2019



**Institute** of  
**mathematics**  
& its applications

# In this talk . . .

- Some facts, well-known and less-well-known, about recurrence relations and power series
- Counting walks on graphs
- Counting words with forbidden factors
- An application



# In this talk . . .

- Some facts, well-known and less-well-known, about recurrence relations and power series
- Counting walks on graphs
- Counting words with forbidden factors
- An application



# In this talk . . .

- Some facts, well-known and less-well-known, about recurrence relations and power series
- Counting walks on graphs
- Counting words with forbidden factors
- An application



# In this talk . . .

- Some facts, well-known and less-well-known, about recurrence relations and power series
- Counting walks on graphs
- Counting words with forbidden factors
- An application



# A curiosity?

$$\frac{1}{89} = 0.011235\dots$$



# A curiosity?

$$\frac{1}{89} = 0.011235955056179\dots$$



# A curiosity?

$$\frac{1}{89} = 0.0112358$$

13  
21  
34  
55  
...





# A curiosity?

Let  $F_n$  denote the  $n$ th Fibonacci number (with  $F_0 = 0$ ).

$$\frac{1}{89} = \sum_{n=0}^{\infty} F_n 10^{-n-1} .$$



# A curiosity?

This is simply the power series expansion

$$\frac{z}{1 - z - z^2} = \sum_{n=0}^{\infty} F_n z^n$$

evaluated at  $z = \frac{1}{10}$ . The occurrence of the polynomial  $1 - z - z^2$  should come as no surprise, as it is the driving polynomial of the Fibonacci sequence.



# Not a curiosity

The power series with coefficients which satisfy a linear recurrence relation are precisely the formal power series attached to rational functions: indeed, the denominator of the rational function will be the driving polynomial of the linear recurrence.



# Not a curiosity

The polynomials over a field  $K$  form a  $K$ -algebra with dual the space of infinite sequences.

The linear recurrent sequences form the *restricted* dual, and hence have a coalgebra structure.

Indeed, the polynomials and the linear recurrent sequences form bialgebras and in two different ways.



# Generating functions

Let  $(r_n)$  be a sequence defined by a recurrence relation

$$r_{n+k} = c_{k-1}r_{n+k-1} + \cdots + c_0r_n$$

for  $n \geq 0$ , with  $c_0 \neq 0$ , subject to given initial conditions on  $r_0, \dots, r_{k-1}$ . Let  $f(x)$  be the auxiliary polynomial

$$f(x) = x^k - (c_{k-1}x^{k-1} + \cdots + c_0) .$$

Let the roots of  $f$  be  $\alpha_1, \dots, \alpha_k$ : we shall assume for simplicity of this exposition that  $f$  has no repeated roots.



# Generating functions

The generating function as a formal power series is equal to a rational function of the form

$$\sum_{n=1}^{\infty} r_n z^n = \frac{p(z)}{\tilde{f}(z)}$$

where  $p$  is a polynomial of degree  $\leq k - 1$ , determined by the initial conditions, and  $\tilde{f}$  is the *reciprocal* polynomial

$$\tilde{f}(z) = \prod_{i=1}^k (1 - \alpha_i z) = (-1/c_0) z^k f(1/z).$$



# Traces

We may also express any such generating function as

$$\sum_{i=1}^k \frac{\beta_i}{1 - \alpha_i z},$$

that is, as a *trace*

$$\text{tr} \left( \frac{\beta}{1 - \alpha z} \right).$$



# Impulse response sequence

The *impulse response sequence* has initial values  $0, \dots, 0, 1$ , so that every sequence is a linear combination of this and its left shifts. The corresponding generating function has initial term  $z^{k-1}$ .

We have

$$\text{tr} \frac{1/f'(\alpha)}{(1 - \alpha z)} = z^{k-1} + O(z^k)$$

as the generating function of the IRS in trace form.





# Walks on graphs

Let  $G = (V, E)$  be a directed graph with possibly multiple edges, and  $A$  the adjacency matrix with entry  $A_{ij}$  counting the number of edges from vertex  $v_i$  to vertex  $v_j$ . Then the entries of the matrix power  $A^n$ , for  $n \geq 0$ , count the number of directed paths of length  $n$  between vertices.



# Walks on graphs

We can regard the formal power series

$$\sum_{n \geq 0} A^n z^n$$

as either a formal power series over the matrix ring or as a matrix whose entries are formal power series.



# Generating function

Let  $G = (V, E)$  be a directed graph with possibly multiple edges, and

It is then easy to see that

$$(I - Az) \left( \sum_{n \geq 0} A^n z^n \right) = I$$

and since  $I - Az$  is invertible, this matrix formal power series is an expression for the inverse  $(I - Az)^{-1}$ .



# Formal power series

We are considering the formal term  $z^n$  as encoding paths of length  $n$ . Let  $B$  be a matrix in which every entry is a polynomial in  $z$  with integer coefficients and with zero constant term: so that  $Az$  is an example of such a matrix. Then the sum

$$\sum_{n \geq 0} B^n$$

converges (formally speaking in the  $z$ -adic topology) to a formal power series in  $z$  over the ring of matrices.



# Counting paths

If we want to count the number of paths of length  $n$  from a given vertex  $i$  to vertex  $j$ , then we write

$$e_i^\top \left( \sum_{n \geq 0} A^n z^n \right) e_j$$

where  $e_j$  is the coordinate vector for the  $i$ -th position.



# Counting paths

We have the generating function

$$e_i^\top (I - Az)^{-1} e_j$$

and expanding  $(I - Az)^{-1}$  we obtain

$$\frac{1}{\det(1 - Az)} e_i^\top \operatorname{adj}(I - Az) e_j$$

which is a rational function with the characteristic polynomial of  $A$  as denominator.



# Forbidden factors

We give as a worked example the problem of counting words over a finite alphabet with forbidden factors: that is, words with no consecutive substring drawn from a finite forbidden set.



# An application

The specific application that motivated us to look at this problem comes from the problem of counting binary words in *prefix normal binary words*: words over the alphabet  $\{0, 1\}$  with the property that no factor has more occurrences of the symbol  $1$  than the prefix of the same length. This problem is addressed, for example, by Burcsi et al. They obtain an upper bound of the form  $2^{n-\lg n+1}$  by observing that a binary word in prefix form for which the first  $k$  symbols are an initial run of  $1^{k-1}$  followed by a  $0$  necessarily has the property that it has no factor of the form  $1^k$ .





# de Bruijn graphs

We wish to enumerate words whose factors of length  $k$  come from a permissible set  $S$ . The *de Bruijn graph* or state machine  $G_S$  has vertex set  $S$ . There is a directed edge from a vertex  $xY \in S$  to a vertex  $YZ \in S$ , where  $x, z$  denote single letters and  $Y$  denotes a string of length  $k - 1$ : following an edge represents the movement one place of a  $k$ -long window into the word. This reduces the problem to one of enumerating walks in the graph  $G_S$ .



# de Bruijn graphs

We wish to enumerate words whose factors of length  $k$  come from a permissible set  $S$ . The *de Bruijn graph* or state machine  $G_S$  has vertex set  $S$ . There is a directed edge from a vertex  $xY \in S$  to a vertex  $YZ \in S$ , where  $x, z$  denote single letters and  $Y$  denotes a string of length  $k - 1$ : following an edge represents the movement one place of a  $k$ -long window into the word. This reduces the problem to one of enumerating walks in the graph  $G_S$ .



# Words without runs

We illustrate the application of power series to obtain this upper bound for binary words with no factor  $\mathbf{1}^k$ .

The de Bruijn graph has  $2^k - 1$  vertices, labelled by all binary strings except  $\mathbf{1}^k$ , and each vertex  $xY$  has two arrows to vertices  $Y\mathbf{0}$  and  $Y\mathbf{1}$ : except for the vertex  $\mathbf{01}^{k-1}$  which has only a single arrow, to  $\mathbf{1}^{k-1}\mathbf{0}$ .



# Words without runs

Fix  $k \geq 2$  and consider the number  $F(n)$  of binary words of length  $n$  with no runs of  $k$  consecutive 1. We know that  $F$  satisfies a linear recurrence relation which we can obtain from its de Bruijn diagram. In this case, it is easy to see that  $F$  satisfies the recurrence

$$F(n) = F(n-1) + \dots + F(n-k)$$

with initial conditions  $F(i) = 2^i$ ,  $i = 0, \dots, k-1$ .



# The characteristic equation

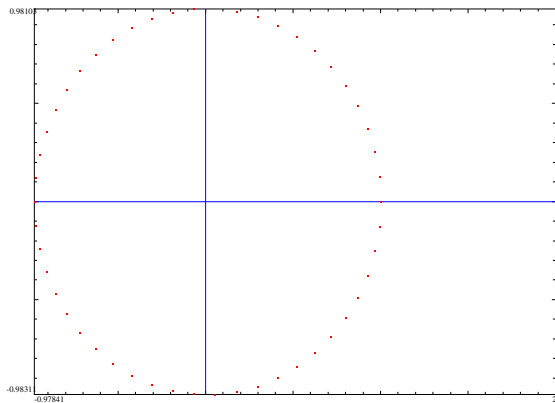
Let  $f(z) = z^k - (z^{k-1} + \dots + 1)$  be the auxiliary polynomial of this recurrence. It is clear that  $f$  has no roots of absolute value greater than or equal to 2. Let  $\alpha_1$  denote the largest real root.

## Theorem

- *The dominant root satisfies  $\alpha_1 < 2 - \frac{1}{2^k}$*
- *$f$  has no repeated roots*
- *All roots other than  $\alpha_1$  are inside the unit circle.*



# A picture



# The characteristic equation

This already gives us the asymptotic behaviour of  $F$ . For a more explicit bound we need to locate the other roots and to estimate  $f'(\alpha)$  as  $\alpha$  runs over the roots of  $f$ .

## Theorem

*The roots of  $f$  within the unit circle are at distance of order  $1/k$  from the  $k - 1$  non-trivial  $k$ -th roots of unity.*



# Conclusion

We conclude that  $f'(\alpha) \sim k$ , for  $\alpha$  one of the roots inside the unit circle. Hence we can obtain an estimate strong enough to give Theorem 15 of [BFLRS]

$$\sum_{k=1}^n F_k(n-k) < 2^{n-\lg n+1},$$





# Summary

- Recurrent sequences and the associated power series have a rich natural algebraic structure
- This structure is useful in a wide range of applications in combinatorics



# Summary

- Recurrent sequences and the associated power series have a rich natural algebraic structure
- This structure is useful in a wide range of applications in combinatorics



# The end

Questions?

Comments!



**Institute** of  
**mathematics**  
& its applications